DIFFERENTIAL PRIVACY: A PRIMER FOR A NON-TECHNICAL AUDIENCE



Kobbi Nissim

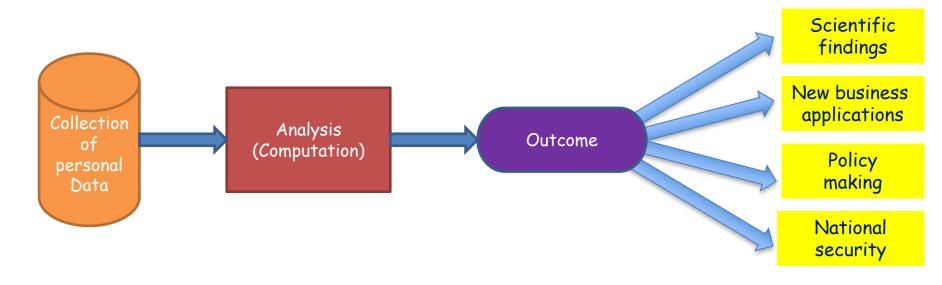
Department of Computer Science Georgetown University



Berkman Klein Center for Internet & Society
Harvard University

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DATA PRIVACY: THE PROBLEM



Given a dataset with sensitive personal information, how can one compute and release functions of the dataset while protecting individual privacy?

ATTACKS ON SDLITECHNIQUES

- Re-identification [Sweeney '00, ...]
 - GIC data, health data, clinical trial data, DNA, Pharmacy data, text data, registry information...
- Blatant non-privacy [Dinur, Nissim '03], 4.2.
- Auditors [Kenthapadi, Mishra, Nissim '05]
- AOL Debacle '06
 - Genome-Wide association studies (GWAS) [Homer et al COS]
- Netflix award [Narayanan, Shmatikov '09]
- Social networks [Backstrom, Dwork, Kleinberg 11]
- Genetic research studies [Gymrek, McGuire, Golan, Halperin, Exhibitation]
- Microtargeted advertising [Korolova 1]
- Recommendation Systems [Calandrino, Kiltzer, Naryapan, Fellen, Shpiatikov 11]
- Israeli CBS [Mukatren, Nissim, Salman, Tromer '14]
- Attack on statistical aggregates [Homer et al Q8] [Work, Smith, Steinke, Vadhan '15]

TAKEAWAYS FROM PRIVACY FAILURES

- Lack of rigor leads to unanticipated privacy failures.
 - New attack modes emerge as research progresses.
 - Redaction of identifiers, release of aggregates, etc. is insufficient.
 - Must take auxiliary information into consideration.
- Any useful analysis of personal data must leak some information about individuals.
- Leakages accumulate with multiple analyses/releases.

Mathematical facts, not matters of policy

ZERO PRIVACY?

Is this where we're headed?



NOT GIVING UP, SCOTT

A new line of privacy work in theoretical computer science (beginning ~2003)

Yields new concept: Differential privacy (2006)

- Rich theory
- In first stages of implementation and real-world use
 - US Census, Google, Apple, Uber, ...

WHAT IS DIFFERENTIAL PRIVACY?

Differential privacy is a definition (i.e., standard) of privacy

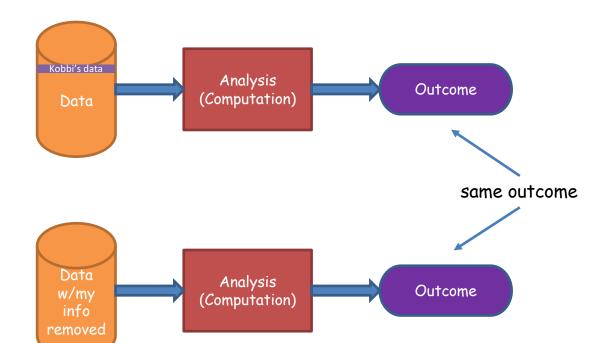
Not a specific technique or algorithm!

Differential privacy is a definition (i.e., standard) of privacy

It expresses a specific desiderata of an analysis:

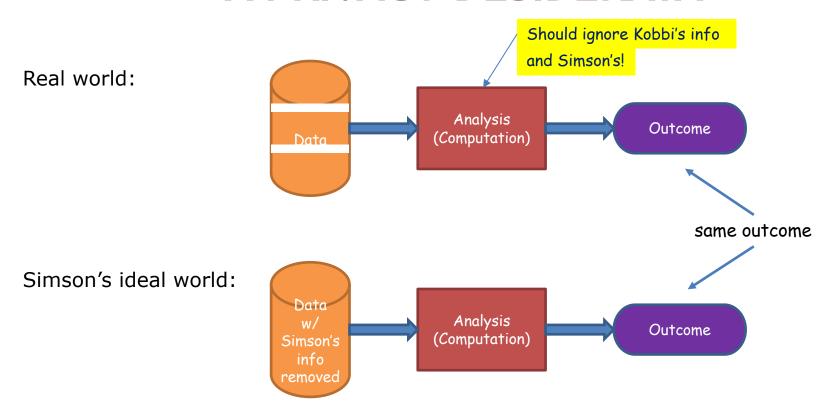
Any information-related risk to a person should not change significantly as a result of that person's information being included, or not, in the analysis.

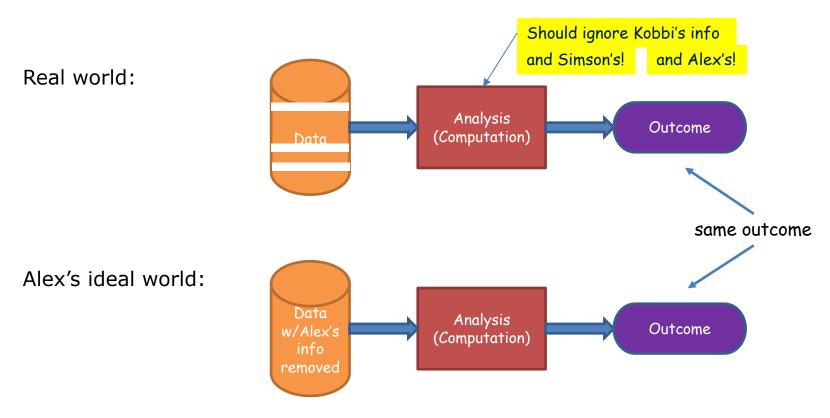
Real world:



My ideal world:

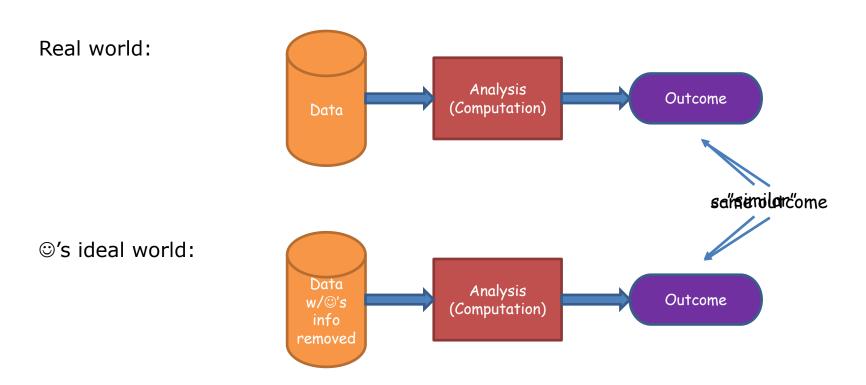
Should ignore Kobbi's info Real world: Analysis Outcome (Computation) Data same outcome My ideal world: Data Analysis Outcome (Computation) info removed



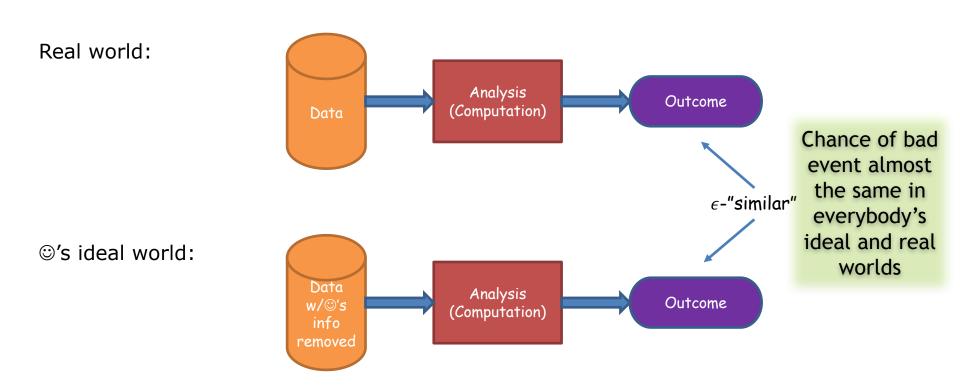


Should ignore Kobbi's info and Simson's! and Alex's! Real world: ... and everybody's! Analysis Outcome (Computation) Data same outcome ©'s ideal world: Data Analysis Outcome w/©'s (Computation) info removed

A MORE REALISTIC PRIVACY DESIDERATA



DIFFERENTIAL PRIVACY [DWORK MCSHERRY NISSIM SMITH '06]



UNDERSTANDING DIFFERENTIAL PRIVACY

- "Automatic" opt-out: I am protected (almost) as if my info is not used at all.
- Plausible deniability: I can claim any value for my information as outcome is (almost) as likely with that value.
- I incur limited risk: Contributing my real info can increase the probability I will be denied insurance by at most 1%.
 - When compared with not participating, or contributing fake info.

- PII: Differential privacy can be interpreted as ensuring that using an individual's data will not reveal (almost) any personally identifiable information that is specific to her.
 - Here, specific refers to information that cannot be inferred unless the individual's information is used in the analysis.

✓ PII

- Linkage: Microdata or contingency tables that allow the identification of population uniques cannot be created using statistics produced by a differentially private tool.
 - This can be formalized and proved mathematically.

- ✓ PII
- ✓ Linkage
- Inference: Differential privacy masks the contribution of any single individual, making it impossible to infer (almost) any information specific to an individual, including whether an individual's information was used at all.

- ✓ PII
- ✓ Linkage
- ✓ Inference

Differential privacy provides protection (far) beyond "identifiability."



Gertrude:

- Age: 65
- She has a \$100,000 life insurance policy.
- She is considering participating in a medical study but is concerned it may affect her insurance premium.



- Based on her age and sex, she has a 1% chance of dying next year. Her life insurance premium is set at $0.01 \times 100,000 = 1,000$.
- Gertrude is a coffee drinker. If the medical study finds that 65-year-old female coffee drinkers have a 2% chance of dying next year, her premium would be set at \$2,000.
 - This would be her baseline risk: Her premium would be set at \$2,000 even if she were not to participate in the study.
- Can Gertrude's premium increase beyond her baseline risk?
 - She is worried that the study may reveal more about her, such as that she *specifically* has a 50% chance of dying next year. This can increase her premium from \$2,000 to \$50,000!



- Reasoning about Gertrude's risk
 - Imagine instead the study is performed using differential privacy with ε = 0.01.
 - The insurance company's estimate of Gertrude's risk of dying in the next year can increase to at most

$$(1+\epsilon)\cdot 2\% = 2.02\%.$$

Her premium would increase to at most \$2,020. Therefore,
 Gertrude's risk would be ≤ \$2020 - \$2000 = \$20.



- Generally, calculating one's baseline is very complex (if possible at all).
 - In particular, in our example the 2% baseline depends on the potential outcome of the study.
 - The baseline may also depend on many other factors Gertrude does not know.
- However, differential privacy provides simultaneous guarantees for every possible baseline value.
 - The guarantee covers not only changes in Gertrude's life insurance premiums, but also her health insurance and more.

COMBINING DIFFERENTIALLY PRIVATE ANALYSES

Combination of ϵ -differentially private computations results in differential privacy (with larger ϵ).

This is extremely important for privacy.

It is a (unique) feature of differential privacy.

Most, if not all, other known definitions of privacy do not measure the cumulative risk from multiple analyses/releases.

THE "PRIVACY BUDGET"

The parameter ϵ measures leakage and can be treated as a "privacy budget" which is consumed as analyses are performed.

Theorems help manage the budget by providing a bound on the overall use of the privacy budget.

This is a feature, not a bug!

Consider how ignoring the fuel gauge would not make your car run indefinitely without refueling.

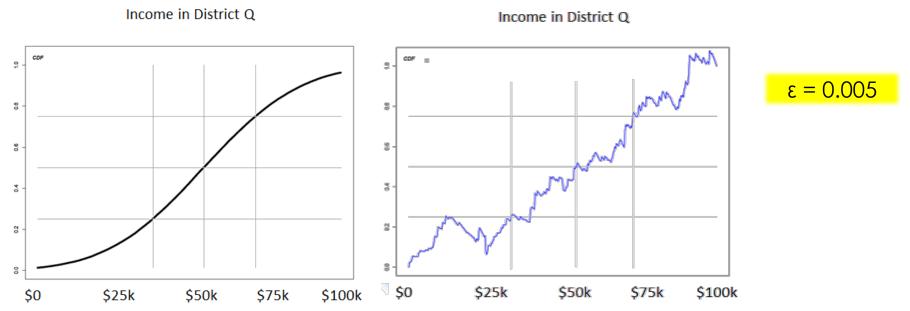


HOW IS DIFFERENTIAL PRIVACY ACHIEVED?

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DIFFERENTIALLY PRIVATE COMPUTATIONS

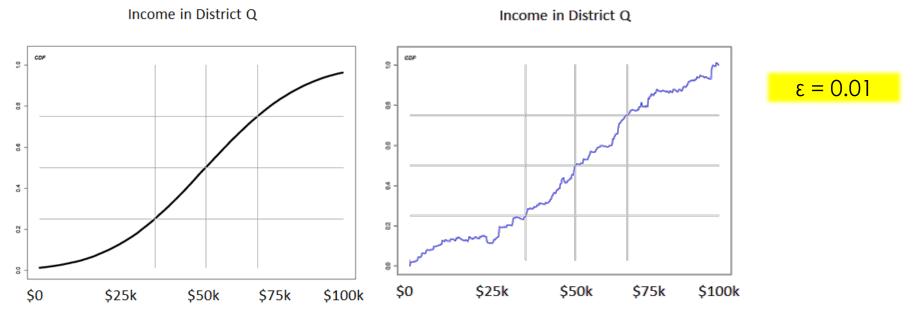
Algorithms maintain differential privacy via the introduction of *carefully crafted* random noise into the computation.



(District Q and its data are stylized examples.)

DIFFERENTIALLY PRIVATE COMPUTATIONS

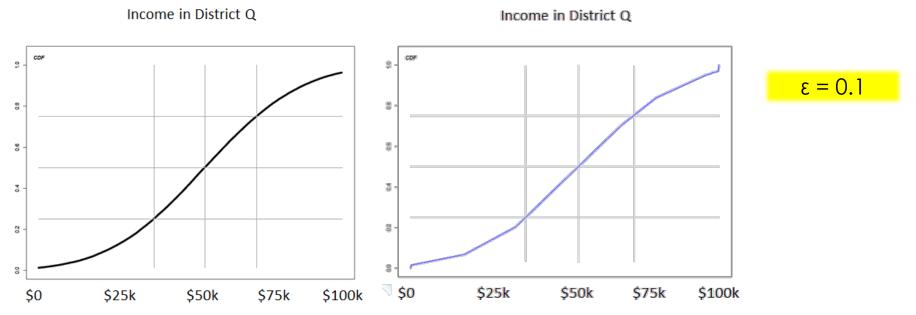
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DIFFERENTIALLY PRIVATE COMPUTATIONS

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WHAT CAN BE COMPUTED WITH DIFFERENTIAL PRIVACY?

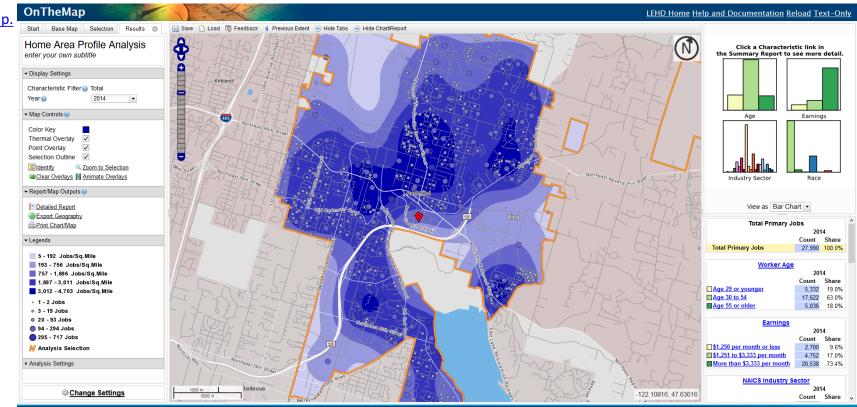
- Descriptive statistics: counts, mean, median, histograms, boxplots, etc.
- Supervised and unsupervised ML tasks: classification, regression, clustering, distribution learning, etc.
- Generation of synthetic data

Because of noise addition, differentially private algorithms work best when the number of data records is large.

EXISTING APPLICATIONS

U.S. CENSUS BUREAU

http://onthemap.ces.census.gov



GOOGLE

RAPPOR: Randomized Aggregatable Privacy-Preserving Ordinal Response

Úlfar Erlingsson Google, Inc. ulfar@google.com Vasyl Pihur Google, Inc. vpihur@google.com

Aleksandra Korolova University of Southern California korolova@usc.edu

ABSTRACT

Randomized Aggregatable Privacy-Preserving Ordinal Response, or RAPPOR, is a technology for crowdsourcing statistics from end-user client software, anonymously, with strong privacy guarantees. In short, RAPPORs allow the forest of client data to be studied, without permitting the possibility of looking at individual trees. By applying randomized response in a novel manner, RAPPOR provides the mechanisms for such collection as well as for efficient, high-utility analysis of the collected data. In particular, RAPPOR permits statistics to be collected on the population of client-side strings with strong privacy guarantees for each client, and without linkability of their reports.

This paper describes and motivates RAPPOR, details its differential-privacy and utility guarantees, discusses its practical deployment and properties in the face of different attack models, and, finally, gives results of its application to both synthetic and real-world data.

1 Introduction

Crowdsourcing data to make better, more informed decisions is becoming increasingly commonplace. For any such crowdsourcing, privacy-preservation mechanisms should be

asked to flip a fair coin, in secret, and answer "Yes" if it comes up heads, but tell the truth otherwise (if the coin comes up tails). Using this procedure, each respondent retains very strong deniability for any "Yes" answers, since such answers are most likely attributable to the coin coming up heads; as a refinement, respondents can also choose the untruthful answer by flipping another coin in secret, and get strong deniability for both "Yes" and "No" answers.

Surveys relying on randomized response enable easy computations of accurate population statistics while preserving the privacy of the individuals. Assuming absolute compliance with the randomization protocol (an assumption that may not hold for human subjects, and can even be nontrivial for algorithmic implementations [23]), it is easy to see that in a case where both "Yes" and "No" answers can be denied (flipping two fair coins), the true number of "Yes" answers can be accurately estimated by 2(Y-0.25), where Y is the proportion of "Yes" responses. In expectation, respondents will provide the true answer 75% of the time, as is easy to see by a case analysis of the two fair coin flips.

Importantly, for one-time collection, the above randomized survey mechanism will protect the privacy of any specific respondent, irrespective of any attacker's prior knowl-

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The Privacy Tools Project is a broad effort to advance a multidisciplinary understanding of data privacy issues and build computational, statistical, legal, and policy tools to help address these issues in a variety of contexts. It is a collaborative effort between Harvard's Center for Research on Computation and

LATEST NEWS & BLOG POSTS

Graduate Student Michael Bar-Sinai Presented at the 8th Annual ESPAnet Israel 2017

PI Salil Vadhan, PI Kobbi Nissim, and Senior Researcher Marco Gaboardi Presented at the Third Biennial Secure and Trustworthy CyberSpace Principal Investigators' Meeting (SaTC PI Meeting '17)

Berkman Klein Center Seeks Applications for 2017 Summer Internship Program

Harvard Magazine Highlights Privacy Tools Project in Article on Privacy and Security

George Kellaris Featured on CRCS Blog

Privacy Tools Project Featured in Harvard Law Review

Berkman Klein Center Seeks Fellow for Privacy

DP IN PRACTICE: CHALLENGES

TRANSITIONING TO PRACTICE

- A new concept:
 - How to communicate its strengths and limitations?
 - What are the "right" use cases for implementation at this stage?
- Access to data:
 - Via a mechanism; Noise added
 - Limited by the "privacy budget"
 - Setting the budget is a policy question
- Matching guarantees with privacy law & regulation

CONCLUSION

MAIN TAKEAWAYS

- Accumulating failures: anonymization & traditional SDL techniques
- Differential privacy:
 - A standard providing a rigorous framework for developing privacy technologies with provable quantifiable guarantees
 - Rich theoretical work, now transitioning to practice
 - First real-world applications and use
 - Not a panacea; to be combined (wisely!) with other technical and policy tools

RESOURCES

LEARNING MORE ABOUT DIFFERENTIAL PRIVACY

- [Nissim et al, 2017] Differential Privacy: A Primer for a Non-technical Audience, Harvard's Privacy Tools project.
- [Dwork 2011] A Firm Foundation for Private Data Analysis, CACM January 2011.
- [Heffetz & Ligett, 2014] Privacy and Data-Based Research, Journal of Economic Perspectives.
- [Dwork & Roth, 2014] The Algorithmic Foundations of Differential Privacy, Now publishers.

+ Online course material, lectures and tutorials.

less technical

technical

PROJECTS, SOFTWARE TOOLS [PARTIAL LIST]

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[Microsoft Research] PINQ
[UT Austin] Airavat: Security & Privacy for MapReduce
[UC Berkeley] GUPT
[CMU-Cornell-PennState] Integrating Statistical and Computational Approaches to Privacy
[US Census] OnTheMap
[Google] Rappor
[UCSD] Integrating Data for Analysis, Anonymization, and Sharing (iDash)
[UPenn] Putting Differential Privacy to Work
[Stanford-Berkeley-Microsoft] Towards Practicing Privacy
[Duke-NISS] Triangle Census Research Network
[Harvard] Privacy Tools
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[Georgetown-Harvard-BU] Formal Privacy Models and Title 13

[Harvard-Georgetown-Buffalo] Computing over Distributed Sensitive Data